Scorpion Maidu: Width Search in the Scorpion Planning System

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This planner abstract describes the Scorpion Maidu (or simply Maidu) planner.\textsuperscript{1} Scorpion Maidu is built on top of the Scorpion planning system (Seipp, Keller, and Helmert 2020), which is an extension of Fast Downward (Helmert 2006). Maidu participated in the satisficing and the agile tracks of IPC 2023. It extends Scorpion with a novel variant of width search algorithms (Lipovetzky and Geffner 2012, 2014, 2017). Maidu uses a different sequential portfolio for each track, combining the new width search algorithms with the configurations used by Fast Downward Stone Soup 2018 (Seipp and Rögner 2018).

Moreover, Maidu replaces the Fast Downward grounder (Helmert 2009) with the new grounder by Corrêa et al. (2023) that uses gringo (Gebser et al. 2011; Kaminski and Schaub 2021).

In this abstract, we only list the new algorithms implemented in Scorpion Maidu, and the settings we used for each track. For a detailed description of the underlying algorithms we refer to the original papers cited below. Scorpion Maidu won the satisficing track of the IPC 2023.

New Width Search Algorithms

Width search (Lipovetzky and Geffner 2012) is based on the concept of novelty: the search is guided towards states that have information not seen before. In classical planning, novelty is computed by tracking which tuples of atoms have been encountered in previous states. The novelty of a state is the size of the smallest tuple of atoms not seen before. As checking and storing all possible tuple sizes is impractical, width-based search algorithms usually track only tuples of size one (single atoms) and two (pairs of atoms).

In general, width search gives a good exploration-exploitation balance while still being cheap to compute. In fact, planners based on width search (Francès et al. 2018) performed very well in the last IPC: LAPKT-DUAL-BFWS achieved second place on the satisficing track, while LAPKT-BFWS-Preference won the agile track.

We introduce new width search algorithms based on forgetting. The rough idea is to forget which tuples have already been achieved from time to time. More precisely, whenever the search makes progress (according to some metric orthogonal to novelty, such as heuristic value or $f$-value), we forget all previously achieved tuples, and start our tracking of tuples from scratch.

Our idea is different from the novelty measures based on partition functions of the search space (Lipovetzky and Geffner 2017; Francès et al. 2017, 2018). Using partition functions, novel tuples are tracked based on the partition (e.g., $f$-value) where they occur. In our algorithm, we do not discriminate the partition where a tuple was first seen, but every time we make progress (in some sense, reach a new partition of the search space) we forget all previous tuples.

We also introduce new ways to implement open lists, trying to create a synergy with our idea of forgetting. It is not useful to forget previous information about novel tuples if the open list is still flooded with very old states. To deal with this, we implemented two different ways to reset the open list once progress is made. First, the more “aggressive” variant simply clears the open list when the search makes progress. Second, in a milder variant, we use a bucket-based queue to implement the open list. Once the search makes progress, the states in the queue are simply moved one bucket back. In this way, new states are inserted ahead of the older ones, even if they have same the heuristic (or novelty) value.

Satisficing Track

The satisficing track has a time limit of 30 minutes, and a memory limit of 8 GiB. In this track, the scores are based not only on whether a plan was found or not, but also on the quality of this plan. To learn our sequential portfolio, we used the Stone Soup algorithm (Helmert et al. 2011; Rögner, Pommerening, and Seipp 2014; Seipp and Rögner 2018). We refer to these planner abstracts and the Stone Soup workshop paper (Helmert, Rögner, and Karpas 2011) for an explanation of how the Stone Soup algorithm works. The learned portfolio uses 20 configurations. It had a total score of 2119.7, while the best single configuration scored only 1574.0. (See the original paper for an interpretation of the scores.)

As mentioned before, we replaced Fast Downward’s grounder (Helmert 2009) used in Scorpion with gringo (Gebser et al. 2011) as done by Corrêa et al. (2023). We also use the $h^2$ preprocessor (Alcázar and Torralba 2015) as a preprocessing step with a time limit of 3 minutes.
Table 1: Coverage comparison between LAMA, Scorpion Maidu, and $h^{new}$, the component configuration of Scorpion Maidu which uses our new method of forgetting fact conjunctions.

<table>
<thead>
<tr>
<th>Task</th>
<th>LAMA</th>
<th>Maidu</th>
<th>$h^{new}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>folding (20)</td>
<td>11</td>
<td>7</td>
<td>10</td>
</tr>
<tr>
<td>labyrinth (20)</td>
<td>1</td>
<td>0</td>
<td>2</td>
</tr>
<tr>
<td>quantum-layout (20)</td>
<td>20</td>
<td>20</td>
<td>20</td>
</tr>
<tr>
<td>recharging-robots (20)</td>
<td>14</td>
<td>14</td>
<td>14</td>
</tr>
<tr>
<td>ricochet-robots (20)</td>
<td>14</td>
<td>17</td>
<td>18</td>
</tr>
<tr>
<td>rubiks-cube (20)</td>
<td>20</td>
<td>20</td>
<td>20</td>
</tr>
<tr>
<td>slitherlink (20)</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>Sum (140)</td>
<td>80</td>
<td>78</td>
<td>84</td>
</tr>
</tbody>
</table>

Figure 1: Cost comparison of best plan found by LAMA with plan found by $h^{new}$.

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Agile Track

The agile track has a time limit of 5 minutes, and a memory limit of 8 GiB. The scores are solely based on how long it takes a planner to find a plan. To learn our sequential portfolio for this track, we used the Greedy Offline Approximation algorithm by Streeter and Smith (2008). Again, we refer to the original paper and the Fast Downward Remix planner abstract (Seipp 2018) for details.

In contrast to the satisficing track, we do not use the $h^2$ preprocessor in the agile track. Furthermore, we also disable the default invariant synthesis (Helmer 2009). However, we still replace the Fast Downward grounder with gringo.

Competition Results

Scorpion Maidu won the satisficing track, so we focus on this track in our discussion. The planner had a total score of 71.86 points, very similar to Levitron’s score of 71.79 (Corrêa et al. 2023a). In fact, Levitron uses Scorpion Maidu as its main component but adds a lifted planner (Corrêa et al. 2023b) as fallback when the grounder fails. Due to this similar performance, both planners were declared joint winners of the satisficing track.

Scorpion Maidu had the highest score among all competitors in the Rubik’s Cube and Slitherlink domains. It was also less than 0.1 points away from the best scoring planners in the Quantum Layout and Recharging Robots domains. On the flip side, it had a score of 0.0 in Labyrinth.

Our portfolio contained only one configuration using the idea of forgetting fact conjunctions as described above. However, this was the most important configuration, as it had an allotted time of almost 10 minutes. We refer to it as $h^{new}$. It is a lazy GBFS alternating between six queues:

1. one queue ordered by the $h^{LMP}$ heuristic (Hoffmann and Nebel 2001);
2. one queue ordered by the $h^{FF}$ heuristic but only containing states generated by preferred operators;
3. one queue ordered by the $h^{LMP}$ heuristic (Richter, Westphal, and Helmert 2011);
4. one queue ordered by the $h^{LMP}$ heuristic but only containing states generated by preferred operators;
5. one type-based queue (Xie et al. 2014) with $h^{FF}$ and $g$-value as types;
6. one queue ordering states by their novelty value (up to width 2) and forgetting previous tuples once progress is made.

To investigate the impact of $h^{new}$, we ran Scorpion Maidu, LAMA, and a planner only running $h^{new}$. We used a time limit of 30 minutes, and a memory limit of 8 GiB, as in the competition. In this experiment, $h^{new}$ also had 30 minutes, instead of the 10 minutes allocated in the portfolio.

Table 1 compares the coverage of the three planners. While the competition had more than one version for several domains (with different PDDL fragments), we used the domains publicly provided by the organizers. Surprisingly, the single configuration $h^{new}$ outperforms both other methods. In three of the seven domains, $h^{new}$ solves strictly more tasks than Maidu.

However, as one might expect, using only $h^{new}$ does not guarantee that we find good plans. Figure 1 compares the cost of the best plan found by LAMA with the (single) plan found by $h^{new}$. It is clear that the higher coverage from $h^{new}$ does not imply good plan quality. A similar situation is observed when comparing the best plan found by the entire Maidu portfolio with $h^{new}$. This indicates that for the satisficing track, $h^{new}$ would not necessarily lead to better scores (as they take into account plan quality) although $h^{new}$ achieves higher coverage.

Acknowledgments

We thank all contributors to the systems that Maidu is built on, mainly Fast Downward (Helmert 2006), gringo (Gebser et al. 2011) and the $h^2$ preprocessor (Alcázar and Torralba 2015).

References

Alcázar, V.; and Torralba, Á. 2015. A Reminder about the Importance of Computing and Exploiting Invariants in Plan-


